Jaipur Engineering College & Research Centre, Jaipur Department of Computer Science & Engineering



Information Security System [6CS4-03] Notes

Prepared By:

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VISION AND MISSION OF INSTITUTE

VISION

To become renowned centre of outcome based learning and work towards academic, professional, cultural and social enrichments of the lives of individual and communities"

MISSION

M1. Focus on evaluation of learning outcomes and motivate students to inculcate research aptitude by project based learning.

M2. Identify areas of focus and provide platform to gain knowledge and solutions based on informed perception of Indian, regional and global needs.

M3. Offer opportunities for interaction between academia and industry.

M4. Develop human potential to its fullest extent so that intellectually capable and imaginatively gifted leaders can emerge in a range of professions.

VISION AND MISSION OF DEPARTMENT

VISION

To become renowned Centre of excellence in computer science and engineering and make competent engineers & professionals with high ethical values prepared for lifelong learning.

MISSION

M1: To impart outcome based education for emerging technologies in the field of computer science and engineering.

M2: To provide opportunities for interaction between academia and industry.

M3: To provide platform for lifelong learning by accepting the change in technologies

M4: To develop aptitude of fulfilling social responsibilities.

COURSE OUTCOMES

On completion of the course, students will be able to:

CO1: Identify different security attacks, Mechanism, classical and modern encryption techniques.

CO2: Apply random number generation, AES and S-box theory and Implement public key cryptosystem.

CO3: Evaluate message authentication and digital signatures using hash function and IP security.

CO4: Analyze & Implement Water marking technique and strong password protocol in Information Security System.

PROGRAM OUTCOMES (PO)

- 1. **Engineering knowledge:** Apply the knowledge of mathematics, science, engineering fundamentals, and an engineering specialization to the solution of complex engineering problems.
- 2. **Problem analysis**: Identify, formulate, research literature, and analyze complex engineering problems reaching substantiated conclusions using first principles of mathematics, natural sciences, and engineering sciences.
- 3. **Design/development of solutions**: Design solutions for complex engineering problems and design system components or processes that meet the specified needs with appropriate consideration for the public health and safety, and the cultural, societal, and environmental considerations.
- 4. **Conduct investigations of complex problems**: Use research-based knowledge and research methods including design of experiments, analysis and interpretation of data, and synthesis of the information to provide valid conclusions.
- 5. **Modern tool usage**: Create, select, and apply appropriate techniques, resources, and modern engineering and IT tools including prediction and modeling to complex engineering activities with an understanding of the limitations.
- 6. **The engineer and society**: Apply reasoning informed by the contextual knowledge to assess societal, health, safety, legal and cultural issues and the consequent responsibilities relevant to the professional engineering practice.
- 7. **Environment and sustainability**: Understand the impact of the professional engineering solutions in societal and environmental contexts, and demonstrate the knowledge of, and need for sustainable development.
- 8. **Ethics**: Apply ethical principles and commit to professional ethics and responsibilities and norms of the engineering practice.
- 9. **Individual and team work**: Function effectively as an individual, and as a member or leader in diverse teams, and in multidisciplinary settings.
- 10. **Communication**: Communicate effectively on complex engineering activities with the engineering community and with society at large, such as, being able to comprehend and write effective reports and design documentation, make effective presentations, and give and receive clear instructions.
- 11. **Project management and finance**: Demonstrate knowledge and understanding of the engineering and management principles and apply these to one's own work, as a member and leader in a team, to manage projects and in multidisciplinary environments.
- 12. **Life-long learning**: Recognize the need for, and have the preparation and ability to engage in independent and life-long learning in the broadest context of technological change.

Program Educational Objectives (PEO)

1. To provide students with the fundamentals of Engineering Sciences with more emphasis in Computer Science & Engineering by way of analyzing and exploiting Engineering challenge

2. To train students with good scientific and engineering knowledge so as to comprehend, analyze, design, and create novel products and solutions for the real life problems.

3. To inculcate professional and ethical attitude, effective communication skills, teamwork skills, multidisciplinary approach, entrepreneurial thinking and an ability to relate engineering issues with social issues.

4. To provide students with an academic environment aware of excellence, leadership, written ethical codes and guidelines, and the self-motivated life-long learning needed for a successful professional career.

5. To prepare students to excel in Industry and Higher education by Educating Students along with High moral values and Knowledge.

Cos/POs	PO1	PO2	PO3	PO4	PO5	PO6	PO7	PO8	PO9	PO10	PO11	PO12
C01	3	3	2	2	1	1	1	1	1	1	1	3
CO2	3	3	3	3	2	1	1	1	1	2	1	3
СО3	3	3	3	3	2	1	1	2	1	2	1	3
CO4	3	3	3	3	2	2	2	2	1	2	1	3

MAPPING CO-PO

Program Specific Outcome's (PSO)

PSO1: Ability to interpret and analyze network specific and cyber security issues, automation in real word environment.

PSO2: Ability to Design and Develop Mobile and Web-based applications under realistic constraints.

Syllabus

SN	Contents	Hours					
1	Introduction: Objective, scope and outcome of the course.	01					
2	Introduction to security attacks: services and mechanism, classical encryption techniques- substitution ciphers and transposition ciphers, cryptanalysis, stream and block ciphers.						
3	Modern block ciphers: Block Cipher structure, Data Encryption standard (DES) with example, strength of DES, Design principles of block cipher, AES with structure, its transformation functions, key expansion, example and implementation. Multiple encryption and triple DES, Electronic Code Book, Cipher Block Chaining Mode, Cipher Feedback mode, Output Feedback mode, Counter mode.	06					
4	Public Key Cryptosystems with Applications: Requirements and Cryptanalysis, RSA cryptosystem, Rabin cryptosystem, Elgamal cryptosystem, Elliptic curve cryptosystem.	06					
5	Cryptographic Hash Functions, their applications: Simple hash functions, its requirements and security, Hash functions based on Cipher Block Chaining, Secure Hash Algorithm (SHA).						
	Message Authentication Codes, its requirements and security, MACs based on Hash Functions, Macs based on Block Ciphers. Digital Signature, its properties, requirements and security, various digital signature schemes (Elgamal and Schnorr), NIST digital Signature algorithm.	05					
6	Key management and distribution: symmetric key distribution using symmetric and asymmetric encryptions, distribution of public keys, X.509 certificates, Public key infrastructure. Remote user authentication with symmetric and asymmetric encryption, Kerberos	04					
	Web Security threats and approaches, SSL architecture and protocol, Transport layer security, HTTPS and SSH.						
	Total	28					

	DEPARTMENT OF COMPUTER SCIENCE ENGINEERING LECTURE PLAN		
Subject:	Information Security System (6CS4-03)	Year/Se	em: III/ VI
Just	No. of Lecture Reqd./(Avl.): 30 / 30	1 001/0	
Semester	Starting: Semester Ending:		
Unit			
No./ Total	Topics to be Delivered	Lect.	Lect. No.
Lecture	Topics to be Denvered	Reqd.	Lett. No.
Reqd.			
Unit-1	Objective, Scope, Outcome of the course.	1	1
(1)	Objective, Scope , Outcome of the course.		
	Introduction to security attacks	1	2
Unit-2 (6)	services and mechanisms	1	3
	Classical encryption techniques	1	4
	substitution ciphers and transposition ciphers,	1	5
	crypt analysis	1	6
	Stream and block ciphers	1	7
	Modern Block Ciphers: Block ciphers structure	1	8
	Data Encryption Standard(DES), Strength of DES	1	9
Unit 3-	Design principle of block cipher	1	10
(6) BC-1	AES with Structure, Key Expansion	1	11
	Multiple Encryption and triple DES	1	12
	Cipher Block Chaining Mode, Cipher feedback mode, Counter mode	1	13
BC-1	IDEA 64 Bit Encryption & MD5 Message Digest Algorithm	1	14
	Public Key Cryptosystems: Requirements	1	15
	Public Key Cryptosystems: Analysis	1	16
Unit 4-	RSA Cryptosystem	1	17
(6)	Rabin Cryptosystem	1	18
	Elgamal Cryptosystem	1	19
	Elliptic Curve Cryptosystem	1	20
	Cryptographic Hash Functions, Hash Function based on Cipher Block Chaining	1	21
Unit 5-	Secure Hash Algorithm	1	22
(5)	Message Authentication Code	1	23
(-)	MAC based on Hash Function & Block Cipher	1	24
	Digital Signature, Various Digital Signature Schemes, NIST Digital Signature	1	25
BC-2	IP Security with Strong Password Protocols	1	26
	Key Management & Distribution, X.509 Certificates	1	27
Unit 6- (4)	Remote User Authentication	1	28
	Web Security Threats, SSL Architecture	1	29
	Transport Layer Security, HTTPs & SSH	1	30
Referenc	es:	· · · ·	
) Stalling	g Williams: Cryptography and Network Security: Principles and Practices, 4th Edit	ion, Pears	on Education
	e & Washington, Introduction to Cryptography, 2nd Ed. Pearson. an Charlie et.al; Network Security: Private Communication in a Public World, 2nd		D

LECTURE PLAN

Unlt-1st ISS - Information Security Attach Security Attack: - An attack is any attempt to destroy, en galu unautuohirzed access. Any action that competendese the Security of information owned by an enganization is repterried to as security attack. cortegenization of atlacks is in term of passève attack & acteue attack. 1> Passive attack 2) Actore Attack 1) <u>Passère</u> attack :- Attacks that do net affect the system are called passène Attack. The teun passene attack indecate that the attacker does net attempt to perferm any modefication to the data. classifuécation of passine attacks into two Sub - contegacées. Passine attack Release of riessage content Traffice Analysis

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" fel (IP) a) Release of Message Contents :- In duis opponent P. Cont They to release on read 1 the contents that its Transnellool. 9/1 feu Example :- A Telephone conveusation is recorded Call ensi er disten by Theed person. Copponent) content C I e en e Sphan Tribernet er opher (Sender) e Ram. e continuestation (Receiver) e facelety e) traffic Analysis :- In this the opponent try C C to analyze and to determine C location & identity of communicating host & C Equila Observe the faceauency and length of message being exchanged. C Car Opponut Sotlan Opponut Cobserve the Patteen of msg) Ram (Receive) Rani (Sendeu) In this type of attacks Actine Attack C. the opponent does the modetton of the data stream ere the 6. creation of anoce False strucim. **S**., 6 classiqueed in four Types. Actine an C.

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Masqueerade :- This active attacks takes place -a) uchen opponent toy to peretend Street the Régénal Senden & The contents of the enjenal message ane modified in some may. four Example. auturd'eateon Servence can be Corptivued and replayed after a valed authentication Servence. (opponent) to be from Solian Internet ou ofner communationeng facelety us facelety man J Rame roleie C (Receiver) 6- It involves the passive capture of a data mult 4 ils subsequent Keplay 65 en unautheuized effect. Capture Teapture message richand [Capture message From to Copponent) [Ram roter Replay message to 2007 message to Ram] nau (sender) Math m Ram (Receiver) Internet en other communicating facelety Modéfication :- let means the some peuteon of aligenal message is alterned, déloyed, récorded to peroduce an unauthoused effect.

modefied the copponent) Say UTH ent I E Internet Sohan C. All Ram (sender) CAR (Receiver) CA C 1 riehan the opponent medéquées the message C Being Tecansgeening 5/10 Selian & Ram. C Cop d) Denial of Sennice :- It prevent en intabet 6 C the neumal die en management of e Communication faceleties. This attack many e e have a spécific tauget Ex:- A message C C in Suppliess in a pauticulair deucction Carry as a securety andet secure. C rechan discupts senico C C Prioriedeal by server C richan (opponent) Server (Sender) Server Internet 2

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Securety Secures of Securety secures nucleur enviores supplement Security of the system en of data being 1 ransfervied Accending to X.800 d'ulder these security Securices indo Fine categories. 1) Authentication :- The authint cation process ensure that the ougen of electroules message et document is connectly identified. it puonde the two specific semice:-(1) <u>Peece entery autred'cation</u> - use in logical Connection to puoulde confidence in the clotentity of the entities connected. is Data eugen autheut cateon :in a connection less Transfer, puorede assured that the source of received data ils as claimed. A I am B. UserA I C-1 Absence of authentication

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011. e/11 Example ?- User à sends au electronile document Over the internet to user B. 4 user e, 1 · AL C has posed himself a user A. UP 4 new pueblem is that how to absen CIL B Know ithat mig bas come fecom ilsur C. CI This type of attack is called fabeleater. CI CO fabrication is possible in absence of COL Puoper authentication mechanisms. CI 2) Access Control :- The puinciple of access o control CI e determines rules should be eible e to access rubert. e e) Data confidenteally :- The puinceple of confidentiality C Specifies that only the sender & the intended recipient (s) should be able Ç to access the contents of a message. confidentiality gets comperoncised if an unauture lised peuson is able to access a meg. c Ex. User A 4 to want to send the mag to C User & but another User c gels access c to tuis msg, much is not descuel, C This defeat the purpose of confidentiality. C Tuis type of attack is called as inter ception. Interception causes loss of message confidentiality C A Sendemoer > B C C C 2088 of Conficol entrally C C C

Data Integerty :- This security semille checks the active attacks. It assures that data received are Exactly as send by an authorized entity. It means that no modepleateon, insenteon, deletion, en replay of the message.

5) Non-repudration? - There are setuation where as a discre Sends a message and dater on Repuise that she had send that mig. Non-Repudration does not allow the sender of a message to repute the Claim of met sending that message.

A flat message, B wulch yeu clouent to name fleceived

Securelty Mechanism: - Some of the Securety mechanism is implemented into the partecular protocol dayer. Security mechanism process called enceptionment. Securety enceptionment. (1) <u>Revensible Enceptionment Mechanism</u>: -<u>A unversible enceptionent Mechanism</u>: -<u>A unversible enceptionent mechanism</u> is simply an enception algeoretum that allows data to be encuypted 4 decrypted.

Coffi in turs worth hash algouething & message 6. authentecation codes, which are used in く digthal signature & message authentecation 6 CA application. C Classical Encouption Techniques :- There are two en en classic / conventional encuyption e techniques of cuyptequaphy. e e 1) Substitution Techniques e 2) Transposition Techniques. e j (i) Substitution Techniques :- A alphanumente chanacter of the plain text is C Replaced by alphanumeulc character. C C (1) Caesar cephen o- Int was due fillest C Example of Subscritteron **C**___ Cephen. The chanacters of plain text C C message aue replaced by efficie character, C sumber de symbols. In this techniques C cach alphabet in a mig is replaced C., by an appracet twee places down C C the eline. LORESH -> ORNHVK Ch (11) Modéquéed veuseon of coursare ceptiere :-In this technique we can not fix the Patterin of Pext. In this we use any same pattern of the text,

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() Mono-Albhabette Cepher :-

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ilt ils simillare to cass Calsare cliphere but in quis cliphere we give any alphabet to this Cliphere.

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Homophonec cepner :-

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in this cephen we can Replace a Plain Text in cephen text, it is abo apphabetic techniques in this we change Bingle apphabet in more than one alphabet for <u>EX</u>. A -> DIH, P, R

 $B \rightarrow E, T, Q, S$ etc.

Pelyguan Substitution Cliphen:-In Alles technique block of alphabet are change in another block of alphabet. EX. HELLO - YUQQW but HELL is retally défenent ceptier text TEUT

Vansposeteon Techniques :- Transposebeon Techniques C.A. ent C Replace the one apprechet nette anether apprechet. in this technique we use different combination of apprechets. 1 C CI C13 e Rall fence Techniques:-C -> In tuis plain Text weate down in en ent Servence of diagonal. em & than seead the plain text in new step (i) -> e e security secur read & weeke it e e e Sequentially $\frac{1}{2} \int_{e}^{m} \int_{o}^{h} \int_{e}^{m} \int_{o}^{m} \int_{o}^{m} \int_{v}^{m} \int_{v}^$ SI e e in ceptier text. CmhmTmrooe O eoorw c c Simple Columnaue Transposeteon rechniques C C (1) Beste Techniques -C - In this we welte the plain text C acous by scow. C → Read ettre message Column by Column 4 ett is net in the endere of 1,2,3 ere 3,2,1, etc. C C C

Oliginal plain text message? - come home Tomercie C. 15

1) Let us conscoler metter 6 columns. To multe the plain text.

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e	se	υ	W	7	

Let il decide the ender of columns as some random ender 4,6,1,2,5,3

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4 than cephen text is cowoocmeroenhometo Simple collimnar transposeteon technique metter mutteple Round

1) wulte a plain text sour by sour in a suctangle of pue-deprine size.

2) met fix the ender of column is 173 or 3,2,1" fou the cephen text.

3) Répeat the (i) & step as many time desiled.

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() <u>versam clephere</u> - This technology is also Called One time pad, 04est CI The length of supert cephere text is canal to the length of plain text. eal 0 Coll O 1 2 3 4 5 6 7 8 9 10 11 12 A B C D E F G H I J K L M Call CI CT ZYX WVUTSRQP D N CI CI 14 13 24 23 22 21 20 19 18 17 16 15 25 C e e Plain Pext H O W A R E YOU e "discourd afdere T 14 22 0 17 4 24 14 20 a single sure at unel ext." One time Pad N C B T Z Q A R X e C **C** c jubbal getal 2016 23 19 [42] 20 24 [3] [43] C c Subtract 26, by 20 16 23 19 16 20 24. 5 17 Vesult >25 U Q X T Q U Y F R C c C C C Cryptanalysis - The perocess of analyze the plain text or Key its Anown as Couptanalysis. C C C C C C 1) Attempt to break a dingle mag C C 2). Recognize tu pattern of encrypted msq, to be able to break (X)

Subsequent ones by applying decuyption of Stream & Block Cloner of The Gremenation of the Cloner text from plain text can be classified in 2 ways. 1) Stream Clepher 2) BLOCK CEPHER 1) Stellary Claphen :- The plain text encrypteal One set at a time. Orejenal meg in plain text is 'lekesh' in ASCIL. 4 New we convect in binaey values, Suppose ilt is 01011100. A key is applied is 10010101 in binary we cassume that we are applying X-OR logte fer encryption. Samelof-0101110.0 Ex. ploe'n text 10010101 key 11001000 2) Black clephine "- A Black of bet is in this plain dext " cakeser _ kinner . searing" f the block copiese cokesh well be encuypted Fast rellowed by know & Finally Shanny

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Déffusion :- This ils also queuen as Transposites when the posteron of letters in the Plaintext of letter in the plaintext X = (Xo, X1, -- Xm-1) au decoverged (Xo, X, ; Xm) -> (Xxo, Xx1, 1. Xxn-1). Diggueston increase the redundancy of the plaintext by spreading ilt access nous & columns. Feistel Cepher steucture :- Most of the block claner used tub claire staucture The Input to the encuyption algorithm is a plaintext block of 2B bets length and key The plain text block is divided into 2 pauts of B bills each as Lo & Ro. These 2 Parts of data are processed through no reminds. After n remain of publicessing the two output block of data and combined to prioduced the cepher text. the cliphere texit. All recends has the same stoucture. The Performance of the felstel depher Atuctuere depender ou the following task. 1) Number of Rounds " The no. of seconds, liquer the secret we use 16 remains.

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et is don't a specific scheme of block appear of Block Size o- Longer block Size underste Security. er plu Y. BET C But it reduce the computational speed STE. of encerption / decuption. Grenerally a block -C 64 bits is used. CIE C Key Bize :- larger me rey size reduced the en encuyption / decuyption speed but it ¢ III Recordes the gueater securety. c Sub-key Greneration Algeuettum - Greceater complexity C in this algerettun lead C to Grueater d'équicanty to cuptanalysis. C 1 C Remol function :- gueater Resistence to the C C Plaintext 2B bits Ro Bblts C LO BEEts key Rend 1 Subkey arneution 1 F K algoer thus C III C C KL C CI C $\bigcirc \oplus$ C C C Bbets Bbets Rn+1 CEL Ln+r-Cepneutex1 (2Bbets) Cold Cold C

Data Encerption Etomolaerol (DES): (Key block) DES is very well known block explice encenption agentum mont published in 1977. by the Matconal Bureau of Standourals for in commenced. and In DES we use the placen text black Moe 64 bets is length and key is of 56 bets in length. longere text is drulded into 64 bet block. 54 bets key V & Feu lauren 64 bet plain text Permited chece 1 Tuetcal Permitation Reind 1 K Pleinuted Chelce 2 Left Cenculde Rend 2 2 Permided chelce 2 k lett Rend 16 Per muted chelce 2 Cencular (Output Bashe Steructure Swap reyet DES & Right hames final permitation 标识 15 PRAME 64 bet cepnen

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4-Tuetial Permutation :- This phase Reaverage the CIL C 64 bet plaintext to peroduce CIE the premited output This is enhance CAR the DES Security Copie S To be overweetten Ble poseteon in the plain en welter the contents Expansion text BLOCK en Key Transfermition this bet Posteron CI Compliess trey Freem Remudators S6 bet to 32 betto RPTfrom en 1 **600** = 58 3264+0 4864 2 C 4864 50 3 42 48 bet -> XOR - 48 ben RPT e Key L Ċ Sbox -Substition 64 ocleet 7 8 7 6 5 4 2 2 2 10 18 26 34 42 58 50 4 12 20 28 36 44 52 60 2 G 14 22 30 38, 46 62 54 3 B 16 56 48 24 40 32 64 4 49 1 57 33 9 41 17 25 5 **G**9 G1 19 43 11 3 35 2T61 \$3 13 21 45 29 S 371 7 23 Ŧ 63 SS 15 31 47 39 8 Justi al permutation Reund Feundation: - The 64-bet permited input passes tweargh 16-

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Itereation. The legit & Right harves of cach 64-blt intermediate value alle Ereated as separate 32-bet quantities. Labeled (1 legd) & R(Right). & Peuferm OXOR.

Tupelt Key (S664) to 64 bel Fevery sta -> 28ber -> D8-1 S6 bet Key 0-1 Teftsor left suft Key slifted -blt Permited RI-1 48-6et Peremuted cuelce 2 Ki fiestel Netweng feud 123456789 1011-15 32-blt RI-L SUD. Of 2 2 Dr Ci 4 ey sted (No. of K-y Blifted bet Apeet Feel Rend) A DES Reund The Beength of DES :and devoded -> Obbeeringthe of DES The -> into & coreas :-1) Key Bize Nature of the algentum Mar -10 use key dêze of 56-bet SEA. Sizep- We That can generate 256 déférent ar Possible keys of the DES Algenettine. The statuce Nature of DES ralgemetrym, us maand secure, 6) Kelailble, & Fleible. A PAKAGE

aubre DES 0- it is very strupple to UT CI Ocidentianal deuble DES. The CI deuble DES does the some thing that CI is DES do Only Onces. Double use 2 Keys, 613 40 Called KI & K2. 40 Contin C Plain Text Encouption Text -> Encurption C Plain Text en K2 1 e Cephere text C Double DES Flust Performes DES on the SI C cleighnal text Using the key KI to get the C Cephere text. It again peuseur on the clipher e e text voing anothere key 1/2. The final output is the caphen text (Encurpted Text) The Deuble DES decuyption is the general Perocess of encryption. Decouption Text Decuption Jourghaf Cephere Floin Text Text Elbenble DES Decemption Teetple DES . In allies DES we use 3 keys. C The plain text fileist encuypted meth C C rey re, then encempted meth 2rd key CI K2, & finally with the key K3. where K, K2, K3 and different to each ether

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fer the Decouption clipher text ruise perf Using key K3, then K2 & finally nelth key Ke 4 Obtain the plain text. DES Encuption text 1 DES to K2 Oliginal Plain Text final Cepnere Rext Cepnere Rext K3 Teleple DES nelty Twee keys Lelple DES nelter Two Keys? -Telepte DES its nigney secure but practically implementation is & difficult. Because ilt requere S6X3 = 168 bet fer key, is displacet. So we Overcome this pieblem Teilple DES nette 2 key. Plain lext Encerption de lext Encerption final Cepner Encuyption Cepner rext